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### Linearized Robust Counterparts of Two-Stage Distribution Problems

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> > July 22th, 2015

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### A CLASSICAL DISTRIBUTION PROBLEM



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### A CLASSICAL DISTRIBUTION PROBLEM

Facility location-transportation model



How can one account for demand uncertainty?

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### ROBUST OPTIMIZATION IS NOW A WELL ESTABLISHED METHODOLOGY



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### A CLASSICAL ROBUST DISTRIBUTION PROBLEM

- ► Robust Facility location-transportation model:
  - $\begin{array}{ll} \max_{I \in \{0,1\}^n, x} & \min_{d \in \mathcal{D}} h(I, x, d) \\ \text{s. t.} & x_i \leq MI_i \ , \ \forall i, \end{array} \quad (FacilitySize \ constraint) \end{array}$
  - where h(I, x, d) is the optimal value of



### STATIC ROBUST OPTIMIZATION

• Consider the following static problem:

$$\max_{x \in \mathcal{X}, y} \quad c^T x + f^T y \tag{1a}$$

s. t. 
$$Ax + By \le Dz \ \forall z \in \mathcal{Z}$$
 (1b)

where we assume *n* decision variables, *J* constraints, and *m* uncertain parameters.

### STATIC ROBUST OPTIMIZATION

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 (1b)

where we assume *n* decision variables, *J* constraints, and *m* uncertain parameters.

If Z := {z ∈ ℝ<sup>m</sup> | Pz ≤ q} is a bounded polyhedral set defined by K constraints, then

Problem (1) 
$$\equiv \max_{x \in \mathcal{X}, y, \Lambda} c^T x + f^T y$$
  
s. t.  $Ax + By + \Lambda q \leq 0$   
 $D + \Lambda P = 0$   
 $\Lambda \geq 0$ ,

where  $\Lambda \in \mathbb{R}^{J \times K}$ .

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### TWO-STAGE ROBUST OPTIMIZATION

• Consider the following two-stage problem:

$$\max_{x \in \mathcal{X}} \min_{z \in \mathcal{Z}} h(x, z)$$
(2)

where

$$h(x,z) := \max_{y} c^{T}x + f^{T}y$$
  
s. t. 
$$Ax + By \le Dz.$$

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### TWO-STAGE ROBUST OPTIMIZATION

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where

$$h(x,z) := \max_{y} c^{T}x + f^{T}y$$
  
s. t. 
$$Ax + By \le Dz$$

• This problem can be represented as

 $\begin{array}{ll} \displaystyle \max_{x \in \mathcal{X}, y(\cdot)} & \displaystyle \min_{z \in \mathcal{Z}} c^T x + f^T y(z) \\ & \text{s. t.} & Ax + By(z) \leq Dz \; \forall z \in \mathcal{Z} \end{array}$ where  $y : \mathbb{R}^m \to \mathbb{R}^{n'}$ 

 
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COMPLEXITY OF TWO-STAGE ROBUST OPTIMIZATION

• Unfortunately, the two-stage robust optimization problem:

$$\begin{array}{ll} \max_{x \in \mathcal{X}, y(\cdot)} & \min_{z \in \mathcal{Z}} c^T x + f^T y(z) \\ \text{s. t.} & Ax + By(z) \leq Dz \; \forall z \in \mathcal{Z} \end{array}$$

is known to be intractable in general (see Ben-Tal et al. 2004 for a proof).

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 Some exact methods have been proposed but without any guarantees about convergence time (see Zeng & Zhao (2013) for a column and constraint generation algorithm) INTRODUCTION LINEARIZED ROBUST COUNTERPART FACILITY LOCATION-TRANSPORTATION INVENTORY MANAGEMENT CON 0000000 000 000 000 000 000

## AFFINELY ADJUSTABLE ROBUST COUNTERPART (BEN-TAL ET AL. 2004)

• Only consider affine recourse functions :

$$y(z) := y + Yz$$

• The two-stage robust problem reduces to

$$\begin{array}{ll}
\max_{x \in \mathcal{X}, y, Y} & \min_{z \in \mathcal{Z}} c^T x + f^T (y + Yz) & (3a) \\
\text{s. t.} & Ax + B(y + Yz) \le Dz \, \forall z \in \mathcal{Z}, & (3b)
\end{array}$$

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(3a)

s. t. 
$$Ax + B(y + Yz) \le Dz \ \forall z \in \mathcal{Z}$$
, (3b)

This model can be reformulated as

$$\max_{x \in \mathcal{X}, y, Y, \lambda, \Lambda} \qquad c^T x + f^T y - q^T \lambda \tag{4a}$$

s. t. 
$$Y^T f + P^T \lambda = 0$$
 (4b)

$$Ax + By + \Lambda q \le 0 \tag{4c}$$

$$D - BY + \Lambda P = 0 \tag{4d}$$

 $\Lambda \ge 0, \lambda \ge 0$  (4e) 32

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### LIFTED AARC (CHEN & ZHANG 2009)

• Consider affine policies on lifted space:

$$y(b) := y + Y^+ z^+ + Y^- z^-$$

where  $z = z^+ - z^-$  and  $(z^+, z^-) \in \mathcal{Z}'$  with

 $\mathcal{Z}' := CvxHull\left(\{(z^+, z^-) \mid \exists z \in \mathcal{Z}, z^+ = \max(0; z), z^- = \max(0; -z)\}\right)$ 

The two-stage robust problem reduces to

$$\max_{x \in \mathcal{X}, y, Y^+, Y^-} c^T x + \min_{(z^+, z^-) \in \mathcal{Z}'} f^T (y + Y^+ z^+ + Y^- z^-)$$
  
s. t.  $Ax + B(y + Y^+ z^+ + Y^- z^-) \le D(z^+ - z^-) \ \forall (z^+, z^-) \in \mathcal{Z}'$ .

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• We omit to present the reformulation...

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### LINEARIZED ROBUST COUNTERPART (LINRC) Let our robust optimization problem take the form

 $\max_{x\in\mathcal{X}} \psi(x) \;,$ 

where

$$\psi(x) := \min_{z \in \mathcal{Z}} \max_{y} \qquad c^{T}x + f^{T}y \qquad (5a)$$
  
s. t.  $Ax + By \le Dz \qquad (5b)$ 

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s. t.  $Ax + By \le Dz \qquad (5b)$ 

$$\psi(x) = \min_{z \in \mathcal{Z}, \lambda \ge 0} \qquad c^T x + z^T D \lambda - (Ax)^T \lambda$$
$$B^T \lambda = f$$

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### LINEARIZED ROBUST COUNTERPART (LINRC)

$$\psi(x) = \min_{\substack{z,\lambda \ge 0 \\ \text{s. t.}}} c^T x + trace(D\lambda z^T) - (Ax)^T \lambda$$
  
s. t.  $Pz \le q$   
 $B^T \lambda = f$ 

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### LINEARIZED ROBUST COUNTERPART (LINRC)

$$\begin{split} \psi(x) &= \min_{\substack{z,\lambda \ge 0 \\ s. t.}} c^T x + trace(D\lambda z^T) - (Ax)^T \lambda \\ s. t. & Pz \le q \\ B^T \lambda = f \\ Pz \lambda^T \le q \lambda^T \\ B^T \lambda z^T = fz^T \\ & \downarrow \\ \psi(x) &\geq \min_{\substack{z,\lambda \ge 0,\Delta \\ s. t.}} c^T x + trace(D\Delta) - (Ax)^T \lambda \\ s. t. & Pz \le q \\ B^T \lambda = f \\ P\Delta^T \le q \lambda^T \\ B^T \Delta = fz^T \end{split}$$

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### LINEARIZED ROBUST COUNTERPART (LINRC)

 The following Linearized Robust Counterpart constitute a tractable conservative approximation to the Two-stage Robust problem:

$$\max_{x \in \mathcal{X}} \min_{\substack{z, \lambda \ge 0, \Delta}} c^T x + trace(D\Delta) - (Ax)^T \lambda$$
  
s. t. 
$$Pz \le q$$
$$B^T \lambda = f$$
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• We omit to provide the reformulation...

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 $B^T \lambda = f$   
 $P\Delta^T \le q\lambda^T$   
 $B^T \Delta = fz^T$ 

- We omit to provide the reformulation...
- ... actually it just so happens that it is exactly equivalent to the AARC reformulation

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# McCormick Envelop Improvement (MCE-LinRC)

• Assuming that  $\mathcal{Z} \subseteq [0,1]^m$  and  $\lambda_i^* \subseteq [0,u_i] \forall i$ , these can help:

$$egin{aligned} &\lambda_i z_j \leq u_i z_j \ orall i, j \ &\lambda_i z_j = \lambda_i - (1-z_j) \lambda_i \geq \lambda_i - (1-z_j) u_i \ orall i, j \end{aligned}$$

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# McCormick Envelop Improvement (MCE-LinRC)

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$$\begin{split} \lambda_i z_j &\leq u_i z_j \; \forall i, j \\ \lambda_i z_j &= \lambda_i - (1 - z_j) \lambda_i \geq \lambda_i - (1 - z_j) u_i \; \forall i, j \end{split}$$

► The problem becomes (MCE-LinRC):

$$\max_{x \in \mathcal{X}} \min_{z, \lambda \ge 0, \Delta} \qquad c^T x + trace(D\Delta) - (Ax)^T \lambda$$
  
s. t. 
$$Pz \le q \ , \ B^T \lambda = f$$
$$P\Delta^T \le q\lambda^T \ , \ B^T \Delta = fz^T$$
$$\Delta \le uz^T$$
$$\Delta \ge \lambda e_m - u(e_m - z)^T$$

• We omit to provide the reformulation...

SDP BASED ROBUST COUNTERPART (SDP-LINRC)

• Assuming that  $\mathcal{Z} \subseteq [0,1]^m$  and  $\lambda_i^* \subseteq [0,u_i] \forall i$ , SDP can help tighten the relaxation

$$\begin{bmatrix} \lambda \lambda^T & \Delta \\ \Delta^T & z z^T \end{bmatrix} \succeq \begin{bmatrix} \lambda \\ z \end{bmatrix} \begin{bmatrix} \lambda^T & z^T \end{bmatrix} \Rightarrow \begin{bmatrix} \Lambda & \Delta & \lambda \\ \Delta^T & Z & z \\ \lambda^T & z^T & 1 \end{bmatrix} \succeq 0$$

### SDP BASED ROBUST COUNTERPART (SDP-LINRC)

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$$\begin{bmatrix} \lambda\lambda^{T} & \Delta \\ \Delta^{T} & zz^{T} \end{bmatrix} \succeq \begin{bmatrix} \lambda \\ z \end{bmatrix} \begin{bmatrix} \lambda^{T} & z^{T} \end{bmatrix} \Rightarrow \begin{bmatrix} \Lambda & \Delta & \lambda \\ \Delta^{T} & Z & z \\ \lambda^{T} & z^{T} & 1 \end{bmatrix} \succeq 0$$
The problem becomes (CDP Lin PC):

► The problem becomes (SDP-LinRC):

$$\max_{x \in \mathcal{X}} \min_{\substack{z,\lambda \ge 0, \Delta, \Lambda, Z}} c^T x + trace(D\Delta) - (Ax)^T \lambda$$
s. t. 
$$Pz \le q \ , \ B^T \lambda = f$$

$$P\Delta^T \le q\lambda^T \ , \ B^T \Delta = fz^T$$

$$\begin{bmatrix} \Lambda & \Delta & \lambda \\ \Delta^T & Z & z \\ \lambda^T & z^T & 1 \end{bmatrix} \succeq 0$$

$$trace(Z) \le z \ , \ trace(\Lambda) \le trace(u\lambda^T + z^T)$$

• We omit to provide the reformulation...

### THINGS TO KNOW ABOUT LINRC

- 1. LinRC is equivalent to AARC.
- 2. If we replace  $z = z^+ z^-$  and  $\mathcal{Z}$  by  $\mathcal{Z}'$ , then LinRC is equivalent to Lifted AARC.
- 3. LinRC can be improved on using linear/conic valid inequalities.
  - ► Ex: MCE-LinRC for facility location-transportation problem
  - ► Ex: SDP-RC in Inventory problem
- 4. Using LinRC, we can find new conditions for optimality of Lifted AARC.
  - Ex: Inventory & newsvendor problems

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## ROBUST FACILITY LOCATION-TRANSPORTATION PROBLEM

$$\begin{array}{ll} \max_{I,x} & \min_{z \in \mathcal{Z}} & h(I,x,z) \\ \text{s. t.} & x_i \leq MI_i, \ \forall i, \ I_i \in \{0,1\} \ \forall i. \end{array}$$

where  $\mathcal{Z}$  captures demand uncertainty and h(I, x, z) is the achieved net profit defined as

$$\begin{split} h(I, x, z) &= \max_{Y} \qquad -(c^{T}x + K^{T}I) + \sum_{i} \sum_{j} (\eta - p_{i} - t_{ij})Y_{ij} \\ \text{s. t.} \qquad \sum_{i} Y_{ij} \leq \bar{d}_{j} + \hat{d}_{j}z_{j}, \ \forall j, \\ \sum_{j} Y_{ij} \leq x_{i}, \ \forall i, \\ Y_{ij} \geq 0, \ \forall i, \forall j. \end{split}$$

### MCE-LINRC VS. LIFTED AARC

100 trials, 10 potential points for facilities, 10 customers.  $\mathcal{Z}(\Gamma) = \{z | \|z\|_{\infty} \le 1, \|z\|_{1} \le \Gamma\}$  where  $\Gamma = 1$  to 10.

	Lifted AARC	MCE-LinRC
A	86%	88%
Max gap	13%	6%
Time (Sec)	3.09	3.72

A: Percentage of trials with optimal solution

### INSPIRATION FOR A PENALTY BASED APPROXIMATION METHOD

 Using MCE-LinRC is equivalent to adding penalized excess variables to uncertain constraint in Lifted AARC.

$$\begin{split} h(I, x, z) &:= \max_{Y, \theta} \qquad -(c^T x + K^T I) + \sum_i \sum_j (\eta - p_i - t_{ij}) Y_{ij} - \sum_j u_j \theta_j \\ \text{s. t.} \qquad \sum_i Y_{ij} \leq \bar{d}_j + \hat{d}_j z_j + \theta_j , \, \forall j \\ \theta_j \geq 0 \,, \, \forall j \,, \end{split}$$

and using a lifted affine decision rule  $\theta_j := S_j^+ z_j^+ + S_j^- z_j^- \forall j$ 

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## MINIMIZING THE MAX OF SUM OF PIECEWISE LINEAR CONVEX FUNCTIONS

Many inventory management problems take the form:

 $\min_{x\in\mathcal{X}} \max_{z\in\mathcal{Z}} h(x,z) \;,$  where h(x,z) is piecewise-linear convex function in both x and z variables.

$$h(x,z) = \sum_{i} \max_{k} (c_{i,k}(x)^{T} z + d_{i,k}(x)) .$$

### MINIMIZING THE MAX OF SUM OF PIECEWISE LINEAR CONVEX FUNCTIONS

Many inventory management problems take the form:

 $\label{eq:hamiltonian} \min_{x\in\mathcal{X}} \ \max_{z\in\mathcal{Z}} \ h(x,z) \ ,$  where h(x,z) is piecewise-linear convex function in both x and z variables.

$$h(x,z) = \sum_{i} \max_{k} (c_{i,k}(x)^{T} z + d_{i,k}(x)) .$$

In fact, one can easily express h(x, z) as a second stage problem

$$\begin{split} h(x,z) &:= \min_{y} \qquad \sum_{i} y_{i} \\ \text{s. t.} \qquad y_{i} \geq c_{i,k}(x)^{T} z + d_{i,k}(x) \ , \forall \, i,k \end{split}$$

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REVISITING LINRC DERIVATIONS WITH  $\Lambda^+ := z^+ \lambda_{i,k}$ 

$$\psi(x) = \max_{\substack{(z^+, z^-) \in \mathcal{Z}', \lambda \ge 0}} \sum_{i,k} \lambda_{i,k} (c_{i,k}(x)^T (z^+ - z^-) + d_{i,k}(x))$$
  
s. t. 
$$\sum_k \lambda_{i,k} = 1, \ \forall i$$

which leads to following when using budgeted uncertainty set

$$egin{aligned} \psi(x) &\leq \max_{\substack{z^+ \geq 0, \, z^- \geq 0, \ \Lambda^+ \geq 0, \, \Lambda^- \geq 0, \, \lambda \geq 0}} &\sum_{i,k} c_{i,k}(x)^T (\Lambda^+_{i,k} - \Lambda^-_{i,k}) + d_{i,k}(x) \lambda_{i,k} \ &\leq n, \, \lambda_{i,k} = 1, \, orall i \ &\leq n, \, \lambda_{i,k} = 1, \, orall i \ &z^+ + z^- \leq 1, \, e_m^T(z^+ + z^-) = \Gamma \ &\sum_k \Lambda^+_{i,k} = z^+, \, \sum_k \Lambda^-_{i,k} = z^-, \, orall i \ &\Lambda^+_{i,k} + \Lambda^-_{i,k} \leq \lambda_{i,k}, \, e_m^T (\Lambda^+_{i,k} + \Lambda^-_{i,k}) = \Gamma \lambda_{i,k}, \, orall i \& A_{i,k} \neq A_{i,k} = 0, \, \lambda_{i,k} = 0, \, \lambda_{i,k$$

### EXAMPLE 1: ROBUST INVENTORY PROBLEM

$$\min_{u} \max_{d \in \mathcal{D}} \sum_{t=1}^{T} c_t u_t + K_t \mathbf{1}_{\{u_t > 0\}} + \max \begin{cases} h_t (x_1 + \sum_{j=1}^{t} (u_j - d_j)), \\ -p_t (x_1 + \sum_{j=1}^{t} (u_j - d_j)) \end{cases}$$

- Objective: Minimizing ordering and shortage/holding cost
- Decision variable (*u<sub>t</sub>*): Stock ordered at the beginning of the *t*th period,
- Uncertainty  $(d_t)$ : Demand during the *t*th period.
- Uncertainty characterization :  $d_t := \overline{d}_t + \widehat{d}_t z$  with  $z \in \mathcal{Z}(\Gamma)$ , the budgeted uncertainty

### SDP-LINRC VS. LIFTED AARC

Comparison of <u>average</u> performance over a set of 1000 randomly generated inventory problem instances with 10 periods.

		Method	
Γ	Interval	Lifted AARC	SDP-LinRC
3	$\leq$ 0.0001% (optimal)	52.6%	56.6%
	$\leq 1\%$ (near optimal)	90.4%	98.1%
	Maximum Gap	4.6%	2.0%
5	$\leq$ 0.0001% (optimal)	57.3%	63.5%
	$\leq 1\%$ (near optimal)	96.6%	99.70%
	Maximum Gap	2.6%	1.3%

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		Method		
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	Maximum Gap	4.6%	2.0%	
5	$\leq 0.0001\%$	57.3%	63.5%	
	≤1%	96.6%	99.70%	
	Maximum Gap	2.6%	1.3%	

### SDP-LINRC VS. LIFTED AARC

Comparison of <u>average</u> performance over a set of 1000 randomly generated inventory problem instances with 10 periods.

		Method		
Γ	Interval	Lifted AARC	SDP-LinRC	
3	$\leq 0.0001\%$	52.6%	56.6%	
	<u>≤</u> 1%	90.4%	98.1%	
	Maximum Gap	4.6%	2.0%	
5	$\leq 0.0001\%$	57.3%	63.5%	
	<u>≤</u> 1%	96.6%	99.70%	
	Maximum Gap	2.6%	1.3%	

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### ROBUST MULTI-ITEM NEWSVENDOR

 The following robust multi-item newsvendor problem can be solved optimally using a LinRC model

$$\max_{x \in \mathcal{X}} \min_{\zeta \in \mathcal{Z}(\Gamma)} \sum_{i} (r_i - c_i) x_i - \max \left\{ \begin{array}{c} (r_i - s_i) (x_i - \bar{d}_i - \hat{d}_i z_i) ,\\ p_i (\bar{d}_i + \hat{d}_i z_i - x_i) \end{array} \right\}$$

when using the budgeted uncertainty set with integer budget.

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when using the budgeted uncertainty set with integer budget.

- Proof relies on fact that vertices of relaxed feasible set for (z<sup>+</sup>, z<sup>-</sup>, Λ<sup>+</sup>, Λ<sup>-</sup>, λ) are integer hence the linearisation Λ<sup>+</sup><sub>i,k</sub> = z<sup>+</sup>λ<sub>i,k</sub> and Λ<sup>-</sup><sub>i,k</sub> = z<sup>-</sup>λ<sub>i,k</sub> are exact due to constraints: λ ∈ [0,1], z<sup>+</sup> ∈ [0,1]<sup>m</sup>, z<sup>-</sup> ∈ [0,1]<sup>m</sup> Λ<sup>+</sup><sub>i,k</sub> ≥ 0, Λ<sup>-</sup><sub>i,k</sub> ≥ 0 ∑<sub>k</sub> Λ<sup>+</sup><sub>i,k</sub> = z<sup>+</sup>, ∑<sub>k</sub> Λ<sup>-</sup><sub>i,k</sub> = z<sup>-</sup>
  - $\Lambda_{i,k}^{+} + \Lambda_{i,k}^{-} \le \lambda_{i,k}$  33 / 37

### DISTRIBUTION FREE MULTI-ITEM NEWSVENDOR

The following distributionally robust problem can also be solved optimally using a LinRC

$$\min_{\mathbf{x}\in\mathcal{X}} \max_{F\in\mathcal{D}} \mathbb{E}_F\left[\sum_i (r_i - c_i)x_i - \max\left\{\begin{array}{c} (r_i - s_i)(x_i - \bar{d}_i - \hat{d}_i z_i), \\ p_i(\bar{d}_i + \hat{d}_i z_i - x_i) \end{array}\right\}\right],$$

where

$$\mathcal{D} = \left\{ F \in \mathcal{M} \middle| \begin{array}{l} \mathbb{P}_{F}(z \in \mathcal{Z}(\Gamma)) = 1 \\ \mathbb{E}_{F}[z] = \mu \\ \mathbb{E}_{F}[(z - \mu)^{+}] \leq \mathbf{r}^{+} \\ \mathbb{E}_{F}[(\mu - z)^{+}] \leq \mathbf{r}^{-} \end{array} \right\}$$

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### Questions & Comments ...

## ... Thank you!