

The Value of Distribution Information in Distributionally Robust Optimization

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CRC in Decision Making under Uncertainty
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Monday, November 9th, 2015

Decision making under uncertainty

Let's consider a decision model that accounts for uncertainty:

$$(SP) \quad \underset{x \in \mathcal{X}}{\text{maximize}} \quad \mathbb{E} [h(x, \xi)]$$

- x is a vector of decision variables in \mathbb{R}^n
- ξ is a vector of uncertain parameters in \mathbb{R}^m
- $h(x, \xi)$ is a profit function

To find an optimal solution, one must develop a stochastic model and solve the associated stochastic program

Difficulties in choosing a distribution model

- Developing an accurate stochastic model requires heavy engineering efforts and might even be impossible
- This motivates the use of a distributionally robust optimization model

$$\text{(DRO)} \quad \underset{x \in \mathcal{X}}{\text{maximize}} \quad \inf_{F \in \mathcal{D}} \mathbb{E}_F[h(x, \xi)] \ .$$

where \mathcal{D} captures exactly what is known of the distribution

Distribution information in data-driven optimization

Many methods have been proposed to convert i.i.d. samples $\{\xi_i\}_{i=1}^M$ into confidence regions for distributions

- Hypothesis testing methods: [(Bertsimas et al., 2015)]

$$\mathcal{S} \ \& \ \{\xi_i\}_{i=1}^M \rightarrow \mathcal{D} := \{F \mid \exists \theta, \psi(F) = \theta, T_\theta(\{\xi_i\}) \leq \gamma(M)\}$$

- Moment based method:

[(Delage and Ye, 2010), (Wiesemann et al., 2014)]

$$\mathcal{S} \ \& \ \{\xi_i\}_{i=1}^M \rightarrow \mathcal{D}_{\text{moment}} := \left\{ F \mid \begin{array}{l} \mathbb{P}(\xi \in \mathcal{S}) = 1 \\ \|\mathbb{E}[\xi] - \hat{\mu}\|_{\hat{\Sigma}^{-1/2}}^2 \leq O\left(\frac{\log(1/\delta)}{M}\right) \\ \mathbb{E}[(\xi - \hat{\mu})(\xi - \hat{\mu})^\top] \preceq \left(1 + O\left(\sqrt{\frac{\log(1/\delta)}{M}}\right)\right) \hat{\Sigma} \end{array} \right\}$$

- Distance/divergence based methods:

[(Ben-Tal et al., 2013), (Mohajerin Esfahani et al. 2015)]

$$\mathcal{S} \ \& \ \{\xi_i\}_{i=1}^M \rightarrow \mathcal{D} := \{F \mid d(F, \hat{F}) \leq \gamma(M)\}$$

Physical ambiguity in Two Urns experiment

Consider that there are two urns in front of you. The two urns contain 100 **BLUE** and **RED** balls in unknown proportions.

Choose among the following three gambles:

- Gamble A: If you draw a **BLUE** ball from urn #1, then you win 180\$, otherwise you win 20\$
- Gamble B: If you draw a **BLUE** ball from urn #1, then you win 200\$, otherwise you win nothing
- Gamble C: If you draw a **BLUE** ball from urn #2, then you win 100\$, otherwise you win nothing

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

Physical ambiguity in Two Urns experiment

Consider that there are two urns in front of you. The two urns contain 100 **BLUE** and **RED** balls in unknown proportions.

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- Gamble C: If you draw a **BLUE** ball from urn #2, then you win 100\$, otherwise you win nothing

Distributionally robust optimization model is:

$$\max_{x \in \{0,1\}^3, x_A + x_B + x_C = 1} \min_{p \in [0,1]^2} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C$$

Value of distribution information

- How can one quantify the value of distribution information ?
 - In Two Urns experiment, what is the value of knowing the proportion of balls in either urn #1 or #2?
 - In data-driven problems, what is the value of acquiring/processing more data?
- This might serve many purposes:
 - Indicate whether it is worth investing in acquisition of additional data
 - Guide the type of data that should be acquired

Outline

- 1 Introduction
- 2 Three Different Measures
- 3 Some Theoretical Properties
- 4 Fleet Mix Optimization
- 5 Conclusion & Future Work

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Three possible measures

Let \mathcal{O} be set of possible information that can be made and $\mathcal{D}(o)$ describe the update rule for the distribution set, such that $\mathcal{D} = \cup_{o \in \mathcal{O}} \mathcal{D}(o)$.

- Worst-case value of information:

$$\text{WC-VDI}(\mathcal{O}) = \min_{o \in \mathcal{O}} \max_{x \in \mathcal{X}} \min_{F \in \mathcal{D}(o)} \mathbb{E}_F[h(x, \xi)] - \max_{x \in \mathcal{X}} \min_{F \in \mathcal{D}} \mathbb{E}_F[h(x, \xi)]$$

- Best-case value of information

$$\text{BC-VDI}(\mathcal{O}) = \max_{o \in \mathcal{O}} \max_{x \in \mathcal{X}} \min_{F \in \mathcal{D}(o)} \mathbb{E}_F[h(x, \xi)] - \max_{x \in \mathcal{X}} \min_{F \in \mathcal{D}} \mathbb{E}_F[h(x, \xi)]$$

- Worst-case regret of not using the information

$$\text{WCR-VDI}(\mathcal{O}) = \max_{o \in \mathcal{O}} \left(\max_{x \in \mathcal{X}} \min_{F \in \mathcal{D}(o)} \mathbb{E}_F[h(x, \xi)] - \min_{F \in \mathcal{D}(o)} \mathbb{E}_F[h(x_0, \xi)] \right)$$

where $x_0 \in \arg \max_{x \in \mathcal{X}} \min_{F \in \mathcal{D}} \mathbb{E}_F[h(x, \xi)]$

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on WC-VDI:

$$\begin{aligned}
 \text{WC-VDI}(\text{Urn}\#1) &= \min_{p_1 \in [0,1]} \max_{x \in \mathcal{X}} \min_{p_2 \in [0,1]} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\
 &\quad - \max_{x \in \mathcal{X}} \min_{p \in [0,1]^2} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\
 &= \min_{p_1 \in [0,1]} \max_{x \in \mathcal{X}} (20 + 160p_1)x_A + (200p_1)x_B - 20 \\
 &= \max_{x \in \mathcal{X}} 20x_A - 20 = 0
 \end{aligned}$$

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on WC-VDI:

$$\text{WC-VDI}(\text{Urn}\#1) = 0$$

$$\begin{aligned} \text{WC-VDI}(\text{Urn}\#2) &= \min_{p_2 \in [0,1]} \max_{x \in \mathcal{X}} \min_{p_1 \in [0,1]} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\ &\quad - \max_{x \in \mathcal{X}} \min_{p \in [0,1]^2} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\ &= \min_{p_2 \in [0,1]} \max_{x \in \mathcal{X}} 20x_A + (100p_2)x_C - 20 = \max_{x \in \mathcal{X}} 20x_A - 20 = 0 \end{aligned}$$

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on WC-VDI:

$$\text{WC-VDI}(\text{Urn}\#1) = 20 - 20 = 0 \quad (\text{i.e. confirm no blue in urn \#1.})$$

$$\text{WC-VDI}(\text{Urn}\#2) = 20 - 20 = 0 \quad (\text{i.e. confirm no blue in urn \#2.})$$

- Conclusion: Distribution information has no value!**

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on BC-VDI:

$$\begin{aligned}
 \text{BC-VDI}(\text{Urn\#1}) &= \max_{p_1 \in [0,1]} \max_{x \in \mathcal{X}} \min_{p_2 \in [0,1]} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\
 &\quad - \max_{x \in \mathcal{X}} \min_{p \in [0,1]^2} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\
 &= \max_{p_1 \in [0,1]} \max_{x \in \mathcal{X}} (20 + 160p_1)x_A + (200p_1)x_B - 20 \\
 &= \max_{x \in \mathcal{X}} 180x_A + 200x_B - 20 = 180
 \end{aligned}$$

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on BC-VDI:

$$\text{BC-VDI}(\text{Urn\#1}) = 180$$

$$\begin{aligned} \text{BC-VDI}(\text{Urn\#2}) &= \max_{p_2 \in [0,1]} \max_{x \in \mathcal{X}} \min_{p_1 \in [0,1]} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\ &\quad - \max_{x \in \mathcal{X}} \min_{p \in [0,1]^2} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\ &= \max_{p_2 \in [0,1]} \max_{x \in \mathcal{X}} 20x_A + (100p_2)x_C - 20 \\ &= \max_{x \in \mathcal{X}} 20x_A + 100x_B - 20 = 80 \end{aligned}$$

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on BC-VDI:

$$\text{BC-VDI}(\text{Urn}\#1) = 200 - 20 = 180 \quad (\text{i.e. confirm all blue in urn \#1.})$$

$$\text{BC-VDI}(\text{Urn}\#2) = 100 - 20 = 80 \quad (\text{i.e. confirm all blue in urn \#2.})$$

- Conclusion: One should count the balls of Urn #1 !**

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on BC-VDI:

$$\text{BC-VDI}(\text{Urn\#1}) = 180 - 20 = 160 \quad (\text{i.e. confirm all blue in urn \#1.})$$

$$\text{BC-VDI}(\text{Urn\#2}) = 100 - 20 = 80 \quad (\text{i.e. confirm all blue in urn \#2.})$$

- Conclusion: One should count the balls of Urn #1 !
- Even when Gamble B is removed as an alternative !

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on WCR-VDI:

$$\begin{aligned}
 \text{WCR-VDI}(\text{Urn\#1}) &= \max_{p_1 \in [0,1]} \max_{x \in \mathcal{X}} \min_{p_2 \in [0,1]} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\
 &\quad - \min_{p_2 \in [0,1]} (20 + 160p_1) \cdot 1 + (200p_1) \cdot 0 + (100p_2) \cdot 0 \\
 &= \max_{p_1 \in [0,1]} \max_{x \in \mathcal{X}} (20 + 160p_1)x_A + (200p_1)x_B - (20 + 160p_1) \\
 &= \max_{p_1 \in [0,1]} \max_{x \in \mathcal{X}} (40p_1 - 20)x_B = \max_{x \in \mathcal{X}} 20x_B = 20
 \end{aligned}$$

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on WCR-VDI:

$$\text{WCR-VDI}(\text{Urn\#1}) = 20$$

$$\begin{aligned} \text{WCR-VDI}(\text{Urn\#2}) &= \max_{p_2 \in [0,1]} \max_{x \in \mathcal{X}} \min_{p_1 \in [0,1]} (20 + 160p_1)x_A + (200p_1)x_B + (100p_2)x_C \\ &\quad - \min_{p_1 \in [0,1]} (20 + 160p_1) \cdot 1 + (200p_1) \cdot 0 + (100p_2) \cdot 0 \\ &= \max_{p_2 \in [0,1]} \max_{x \in \mathcal{X}} (20)x_A + (100p_2)x_B - 20 = \max_{x \in \mathcal{X}} 80x_B = 80 \end{aligned}$$

Value of distribution information in Two Urns exp.

Gamble A		Urn #2	
		Blue	Red
Urn #1	Blue	180	180
	Red	20	20

Gamble B		Urn #2	
		Blue	Red
Urn #1	Blue	200	200
	Red	0	0

Gamble C		Urn #2	
		Blue	Red
Urn #1	Blue	100	0
	Red	100	0

If we could count the balls of one urn, which one should it be ?

- Based on WCR-VDI:

$$\text{WCR-VDI}(\text{Urn}\#1) = 200 - 180 = 20 \quad (\text{i.e. confirm all blue in urn \#1.})$$

$$\text{WCR-VDI}(\text{Urn}\#2) = 100 - 20 = 80 \quad (\text{i.e. confirm all blue in urn \#2.})$$

- Conclusion: One should count the balls of Urn #2 !**

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An ordering of VDI measures

In Two Urns experiment, we noticed that

$$\underbrace{\text{WC-VDI(Urn\#1)}}_0 \leq \underbrace{\text{WCR-VDI(Urn\#1)}}_{20} \leq \underbrace{\text{BC-VDI(Urn\#1)}}_{180}$$

$$\underbrace{\text{WC-VDI(Urn\#2)}}_0 \leq \underbrace{\text{WCR-VDI(Urn\#2)}}_{80} \leq \underbrace{\text{BC-VDI(Urn\#2)}}_{80}$$

Lemma

It is generally the case that

$$\text{WC-VDI}(\mathcal{O}) \leq \text{WCR-VDI}(\mathcal{O}) \leq \text{BC-VDI}(\mathcal{O}) .$$

An ordering of VDI measures

Lemma

It is generally the case that

$$WC\text{-VDI}(\mathcal{O}) \leq WCR\text{-VDI}(\mathcal{O}) \leq BC\text{-VDI}(\mathcal{O}) .$$

Proof:

$$\begin{aligned} WC\text{-VDI}(\mathcal{O}) &= \min_{o_1 \in \mathcal{O}} \max_{x_1} \min_{F \in \mathcal{D}(o_1)} \mathbb{E}_F[h(x_1, \xi)] - \max_{x_2 \in \mathcal{X}} \min_{F \in \mathcal{D}} \mathbb{E}_F[h(x_2, \xi)] \\ &= \min_{o_1 \in \mathcal{O}} \max_{x_1 \in \mathcal{X}} \min_{F \in \mathcal{D}(o_1)} \mathbb{E}_F[h(x_1, \xi)] - \min_{o_2 \in \mathcal{O}} \min_{F \in \mathcal{D}(o_2)} \mathbb{E}_F[h(x_0, \xi)] \\ &= \max_{o_2 \in \mathcal{O}} \min_{o_1 \in \mathcal{O}} \max_{x_1 \in \mathcal{X}} \min_{F \in \mathcal{D}(o_1)} \mathbb{E}_F[h(x_1, \xi)] - \min_{F \in \mathcal{D}(o_2)} \mathbb{E}_F[h(x_0, \xi)] \\ &\leq \max_{o_1 = o_2 \in \mathcal{O}} \max_{x_1 \in \mathcal{X}} \min_{F \in \mathcal{D}(o_1)} \mathbb{E}_F[h(x_1, \xi)] - \min_{F \in \mathcal{D}(o_2)} \mathbb{E}_F[h(x_0, \xi)] \\ &= WCR\text{-VDI}(\mathcal{O}) \end{aligned}$$

An ordering of VDI measures

Lemma

It is generally the case that

$$WC\text{-VDI}(\mathcal{O}) \leq WCR\text{-VDI}(\mathcal{O}) \leq BC\text{-VDI}(\mathcal{O}) .$$

Proof:

$$\begin{aligned} WCR\text{-VDI}(\mathcal{O}) &= \max_{o \in \mathcal{O}} \{ \max_{x_1 \in \mathcal{X}} \min_{F \in \mathcal{D}(o)} \mathbb{E}_F[h(x_1, \xi)] - \min_{F \in \mathcal{D}(o)} \mathbb{E}_F[h(x_0, \xi)] \} \\ &\leq \max_{o_1 \in \mathcal{O}} \max_{o_2 \in \mathcal{O}} \max_{x_1 \in \mathcal{X}} \min_{F \in \mathcal{D}(o_1)} \mathbb{E}_F[h(x_1, \xi)] - \min_{F \in \mathcal{D}(o_2)} \mathbb{E}_F[h(x_0, \xi)] \\ &= \max_{o_1 \in \mathcal{O}} \max_{x_1 \in \mathcal{X}} \min_{F \in \mathcal{D}(o_1)} \mathbb{E}_F[h(x_1, \xi)] - \min_{F \in \mathcal{D}} \mathbb{E}_F[h(x_0, \xi)] \\ &= BC\text{-VDI}(\mathcal{O}) \end{aligned}$$

Many situations where $\text{VDI} = 0$ in the worst case

Lemma

If the feasible set \mathcal{X} is convex and compact and the profit function $h(x, \xi)$ is concave in x , then $\text{WC-VDI}(\mathcal{O}) = 0$.

Proof: Based on Sion's minimax theorem we have that

$$\begin{aligned} \text{WC-VDI}(\mathcal{O}) &= \min_{o \in \mathcal{O}} \max_{x_1 \in \mathcal{X}} \min_{F \in \mathcal{D}(o)} \mathbb{E}_F[h(x_1, \xi)] - \max_{x_2 \in \mathcal{X}} \min_{F \in \mathcal{D}} \mathbb{E}_F[h(x_2, \xi)] \\ &\leq \min_{o \in \mathcal{O}} \min_{F \in \mathcal{D}(o)} \max_{x_1 \in \mathcal{X}} \mathbb{E}_F[h(x_1, \xi)] - \max_{x_2 \in \mathcal{X}} \min_{F \in \mathcal{D}} \mathbb{E}_F[h(x_2, \xi)] \\ &= \min_{F \in \mathcal{D}} \max_{x_1 \in \mathcal{X}} \mathbb{E}_F[h(x_1, \xi)] - \max_{x_2 \in \mathcal{X}} \min_{F \in \mathcal{D}} \mathbb{E}_F[h(x_2, \xi)] = 0. \end{aligned}$$

Some situations where $\text{VDI} = 0$ in the best case

Theorem (Delage et al., 2014)

Let the profit function $h(x, \xi)$ be convex in ξ , and let the distribution set be $\mathcal{D} := \{F \mid \mathbb{E}_F[\xi] = \mu\}$. Then for any information sets of type $\mathcal{O} := \{\gamma \in \mathbb{R}^+ \mid \mathbb{E}_F[\psi(\xi)] \leq \gamma\}$ where $\psi(\cdot)$ is a convex function and $\mathcal{D}(o) \neq \emptyset$ for all $o \in \mathcal{O}$, the value is zero even in the best case.

Proof:

$$\begin{aligned} \text{BC-VDI} &= \max_{o \in \mathcal{O}} \max_{x \in \mathcal{X}} \min_{F \in \mathcal{D}(o)} \mathbb{E}_F[h(x, \xi)] - \max_{x \in \mathcal{X}} \min_{F \in \mathcal{D}} \mathbb{E}_F[h(x, \xi)] \\ &= \max_{o \in \mathcal{O}} \max_{x \in \mathcal{X}} \mathbb{E}_F[h(x, \mu)] - \max_{x \in \mathcal{X}} \mathbb{E}_F[h(x, \mu)] = 0, \end{aligned}$$

since Jensen's inequality ensures that δ_μ (i.e. the Dirac measure centred at μ) always achieves a lower profit than any $F \in \mathcal{D}$, and since this Dirac measure remains feasible when imposing that $\mathbb{E}_F[\psi(\xi)] \leq \gamma$.

Evaluating worst-case regret is NP-hard

Theorem (Delage et al., 2014)

Evaluating $BC\text{-VDI}(\mathcal{O})$ or $WCR\text{-VDI}(\mathcal{O})$ exactly is NP-hard even when $h(x, \xi)$ is concave in x and convex in ξ , and $\mathcal{D}_{\text{moment}}$ is used.

Sketch of proof:

- When the distribution information is perfect,

$$\begin{aligned} WCR\text{-VDI}(\mathcal{O}) &= \max_{F \in \mathcal{D}} \max_{x \in \mathcal{X}} \mathbb{E}_F[h(x, \xi)] - \mathbb{E}_F[h(x_0, \xi)] \\ &= \max_{x \in \mathcal{X}} \max_{F \in \mathcal{D}} \mathbb{E}_F[h(x, \xi)] - \mathbb{E}_F[h(x_0, \xi)]. \end{aligned}$$

- Evaluating $\max_{F \in \mathcal{D}_{\text{moment}}} \mathbb{E}_F[h(x, \xi)]$ is NP-hard for

$$\begin{aligned} h(x, \xi) &:= \max_{y \in \mathbb{R}^m} c^T x + \xi^T y \\ \text{s.t.} \quad &|y_i| \leq x, \quad \forall y \in \{1, 2, \dots, m\} \\ &a^T y = 0. \end{aligned}$$

Tractable bound for worst-case regret

Theorem (Delage et al., 2014)

If the following conditions apply:

- 1 $\mathcal{D}_{\text{moment}}$ is used with $\mathcal{S} \subseteq \{\xi \mid \|\xi\|_1 \leq \rho\}$ and $\|\mathbb{E}_F[\xi] - \hat{\mu}\|_{\hat{\Sigma}^{-1/2}}^2 \leq \gamma_1$
- 2 $h(x, \xi)$ captures a two-stage linear program with cost uncertainty, i.e., $h(x, \xi) := \max_{y \in \mathcal{Y}(x)} c^T x + \xi^T C y$.

then an upper bound for $\text{WCR-VDI}(\mathcal{O})$ can be evaluated

$$\begin{aligned} \text{WCR-VDI}(\mathcal{O}) \leq \min_{s \in \mathbb{R}, q \in \mathbb{R}^m} \quad & s + \hat{\mu}^T q + \sqrt{\gamma_1} \|\hat{\Sigma}^{1/2} q\| \\ \text{s.t.} \quad & s \geq \alpha(\rho e_i) - \rho e_i^T q, \quad \forall i \in \{1, \dots, m\} \\ & s \geq \alpha(-\rho e_i) + \rho e_i^T q, \quad \forall i \in \{1, \dots, m\}, \end{aligned}$$

where $\alpha(\xi) = \max_{x \in \mathcal{X}} h(x, \xi) - (c^T \bar{y}_0 + \xi^T C \bar{y}_0)$ for any $\bar{y}_0 \in \mathcal{Y}(x_0)$.

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Value of distribution info for an airline company

- Fleet mix optimization is a difficult decision problem:
 - Fleet contracts are signed 10 to 20 years ahead of schedule.
 - Many factors are still unknown at that time:
passenger demand, fuel prices, etc.
- Yet, many airline companies sign these contracts based on a single scenario of what the future may be.
- We first show that using the mean value of future profits as a scenario leads to the same solution as DRO with $\mathcal{D}_{\text{moment}}$ with known first moment
- Can we do better by acquiring more information about the distribution ?

Mathematical formulation for fleet mix optimization

The fleet composition problem is a stochastic mixed integer LP

$$\text{Fleet mix} \longrightarrow \underset{x}{\text{maximize}} \mathbb{E} \left[- \underbrace{o^\top x}_{\text{ownership cost}} + \underbrace{h(x, \tilde{p}, \tilde{c}, \tilde{L})}_{\text{future profits}} \right],$$

with $h(x, \tilde{p}, \tilde{c}, \tilde{L}) :=$

$$\begin{aligned} \max_{z \geq 0, y \geq 0, w} \quad & \sum_k \left(\sum_i \underbrace{\tilde{p}_i^k w_i^k}_{\text{flight profit}} - \underbrace{\tilde{c}_k (z_k - x_k)^+}_{\text{rental cost}} + \underbrace{\tilde{L}_k (x_k - z_k)^+}_{\text{lease revenue}} \right) \\ \text{s.t.} \quad & \left. \begin{aligned} w_i^k \in \{0, 1\}, \forall k, \forall i \quad \& \quad \sum_k w_i^k = 1, \forall i \end{aligned} \right\} \text{Cover} \\ & \left. \begin{aligned} y_{g \in \text{in}(v)}^k + \sum_{i \in \text{arr}(v)} w_i^k = y_{g \in \text{out}(v)}^k + \sum_{i \in \text{dep}(v)} w_i^k, \forall k, \forall v \end{aligned} \right\} \text{Balance} \\ & \left. \begin{aligned} z_k = \sum_{v \in \{v | \text{time}(v)=0\}} (y_{g \in \text{in}(v)}^k + \sum_{i \in \text{arr}(v)} w_i^k), \forall k \end{aligned} \right\} \text{Count} \end{aligned}$$

Experiments in fleet mix optimization

We experimented with three test cases :

- ① 3 types of aircrafts, 84 flights, $\sigma_{\tilde{p}_i}/\mu_{\tilde{p}_i} \in [4\%, 53\%]$
- ② 4 types of aircrafts, 240 flights, $\sigma_{\tilde{p}_i}/\mu_{\tilde{p}_i} \in [2\%, 20\%]$
- ③ 13 types of aircrafts, 535 flights, $\sigma_{\tilde{p}_i}/\mu_{\tilde{p}_i} \in [2\%, 58\%]$

Results:

Test cases	WCR-VDI(o)
#1	$\leq 6\%$
#2	$\leq 1\%$
#3	$\leq 7\%$

Conclusions:

- It's wasteful in these problems to invest more than 7% of profits in acquisition of distribution information

Outline

- 1 Introduction
- 2 Three Different Measures
- 3 Some Theoretical Properties
- 4 Fleet Mix Optimization
- 5 Conclusion & Future Work**

Conclusion & future work

- Need for tools that can estimate the value of distribution information
 - The most natural tools are computational intractable
 - Tractable upper bounds for value of perfect distribution information might be available and informative (e.g. fleet-mix optimization)
- Future work:
 - Develop tighter bounds for $WCR-VDI(\mathcal{O})$ with perfect distribution information under $\mathcal{D}_{\text{moment}}$
 - Derive bounds for other distribution sets
 - Design simple procedures for characterizing \mathcal{O} and $\mathcal{D}(o)$ and bounding $WCR-VDI(\mathcal{O})$ in data-driven problem where information consists of samples

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Questions & Comments ...

... Thank you!